

**E.D.P. 2003, Monterey, CA.**



# **Interoperability, Datamodels, and Databases**

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# Glossary

- **Data Model**
  - The native in-core data structures used by an algorithm
- **Database**
  - Client/server disk or memory based data repository
  - (non-native)
- **Application Programming Interface (API)**
  - How software modules interact with data model or database
- **Standard File Format**
  - ASCII or Binary file representation of a data model or database

# Why all the fuss about data storage?

- There is no such thing as a best-in-class point tool, only a best-in-class flow
- “It’s the flow, stupid!”
  - (My apologies to Bill Clinton...)
- EDA = Flow + Data model + Algorithms
  - (If you squint)

# Nvidia Example: Adding more engineers to deal with complexity

Design Start	Technology node	Transistor count	Complexity	Front-end staff	Back-end staff
1993	0.5 $\mu$	0.75M	1x	1.0x	1.0x
1995	0.5 $\mu$	1.25M	1.5x	1.2x	3.0x
1996	0.35 $\mu$	4.0M	4x	1.6x	3.0x
1997	0.31 $\mu$	7.5M	7x	1.7x	4.0x
1998	0.25 $\mu$	9.0M	10x	1.5x	4.0x
1998	0.22 $\mu$	22M	20x	2.5x	5.0x
1999	0.18 $\mu$	25M	22x	1.5x	4.0x
1999	0.15 $\mu$	57M	30x	3.5x	6.0x
2000	0.15 $\mu$	60M	35x	1.5x	7.0x
2000	0.15 $\mu$	63M	40x	3.0x	7.0x
2001	0.13 $\mu$	120M	50x	5.0x	9.0x

Source: Nvidia/Chris Malachovski DAC 2002



# Cannot Justify Traditional Flows for Today's Designs

- **40 M gate designs**
  - 18mm X 18mm, 2000 I/Os, 500Mhz
  - Approximately 4M lines of RTL
- **50+ engineers**
  - Experts in synthesis, P&R, signal integrity, power analysis, design closure
- **\$80MM investment**
  - Requires \$160MM in 2 years to realize break even
  - Where is the killer application for this??

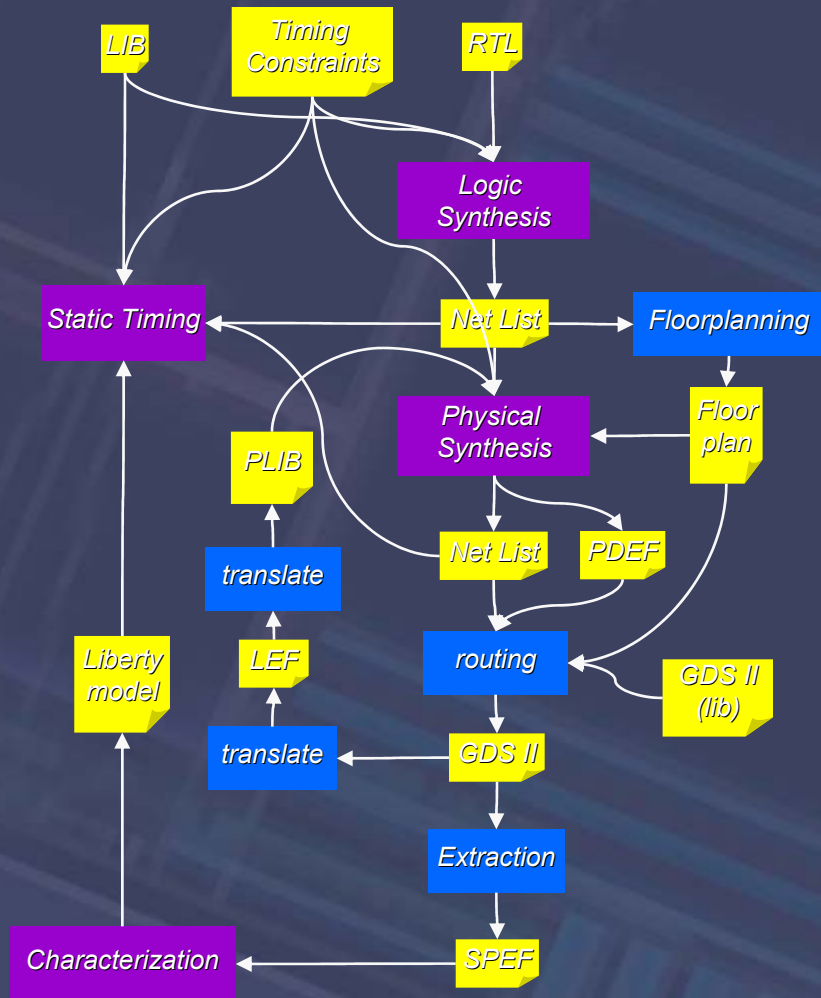
***Traditional design flows will make Moore's law economically infeasible***



# Point Tools = Poor Productivity, Lower QoR

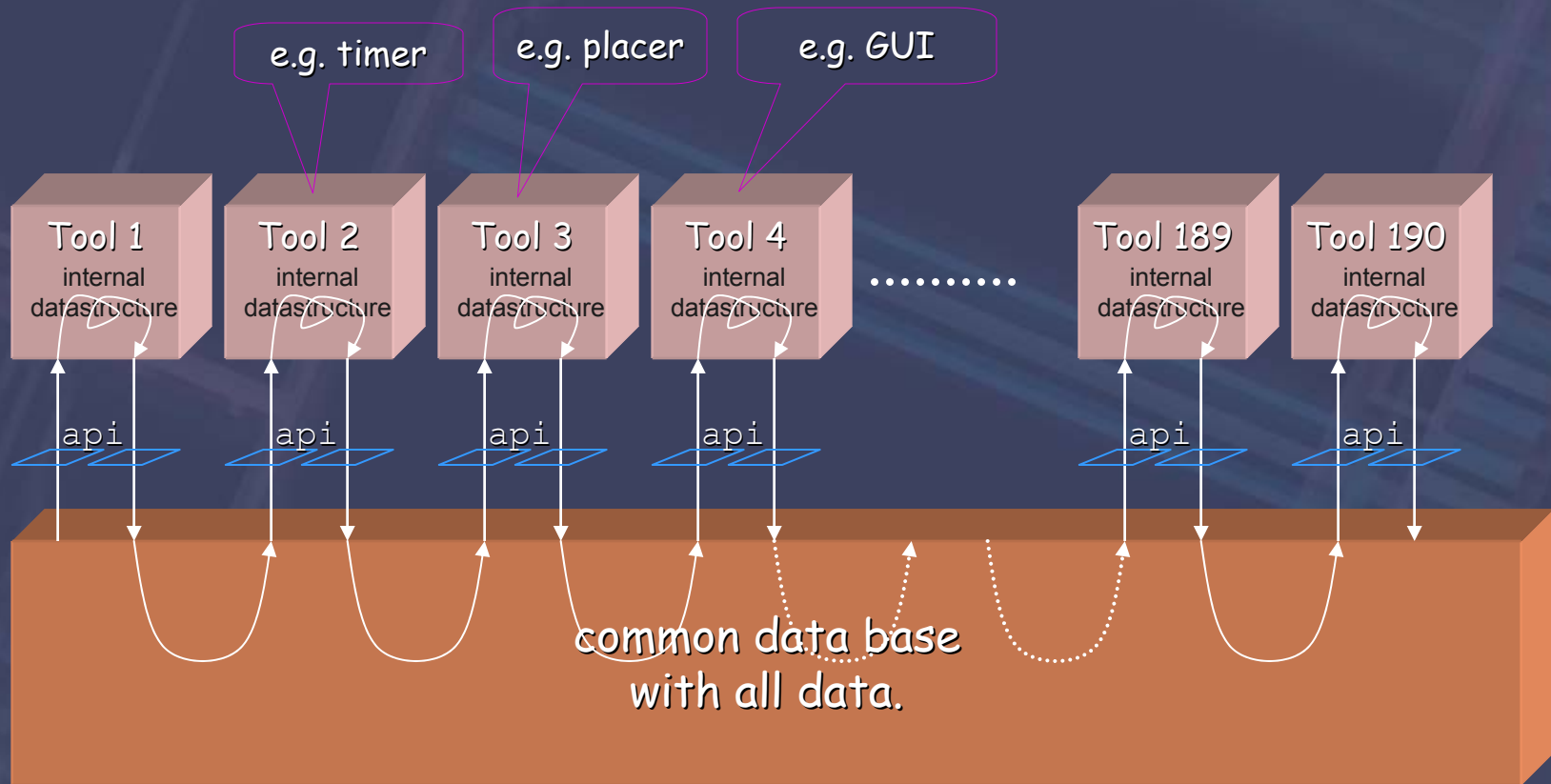
- **Point tools – architecturally inefficient**
  - Not scalable for nanometer effects: each new problem needs to be addressed by a new point tool
  - Requires wasteful iterations
- **Different timers, placers, libraries, constraint systems**
  - Correlation problems
  - Guardband to gain timing predictability → lower quality of results (QoR)
- **Huge file transfers cause waste of time**
  - Several hours of idle time for every iteration

# 'Bolting' together a flow using files



- Design data is spread over many files
  - Files are big and slow to read
- Relevant information gets lost in the translation
  - interpretation may not be consistent
- DSM issues are everywhere
  - the file interface makes dealing with them harder

# Tool integration through a 'Common Database'



- The database interfaces with the tools through an API
- Each tool makes its own copy of the data (data model)
- It takes memory and time to haul all this data around



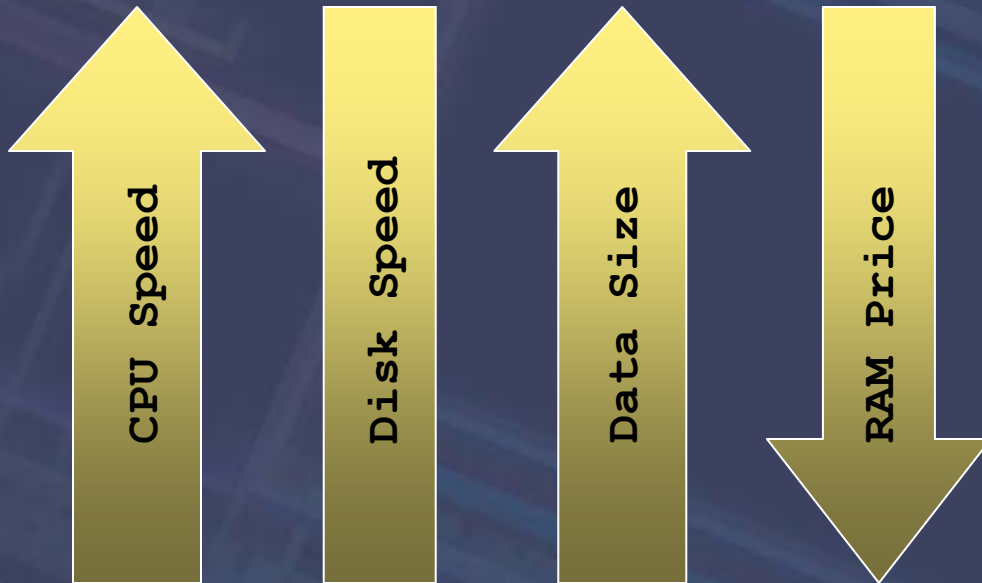
# Would you architect it like that if you could start from scratch?

# NO WAY!

## instead:

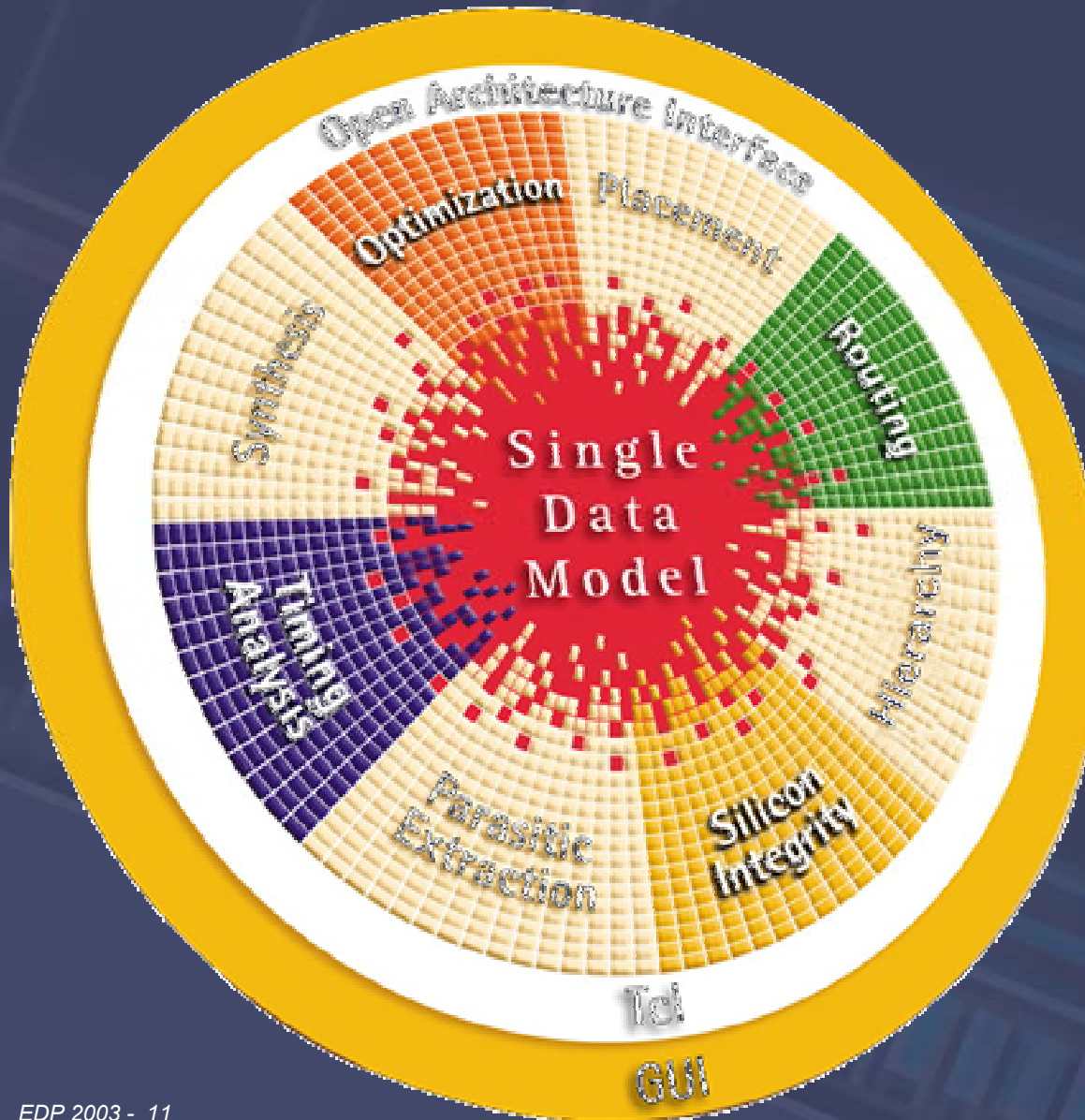
- **Leverage similarity**
  - save implementation effort
  - reduce bugs
  - have consistency by construction
- **Minimize interfaces**
  - Tools spend a LOT of code reading, writing, and conditioning data.
- **Add incremental analysis tools (Timer, extraction, DRC) as part of the infrastructure. They are not point tools, they are part of the datamodel!**

# Why is it time for a paradigm shift?



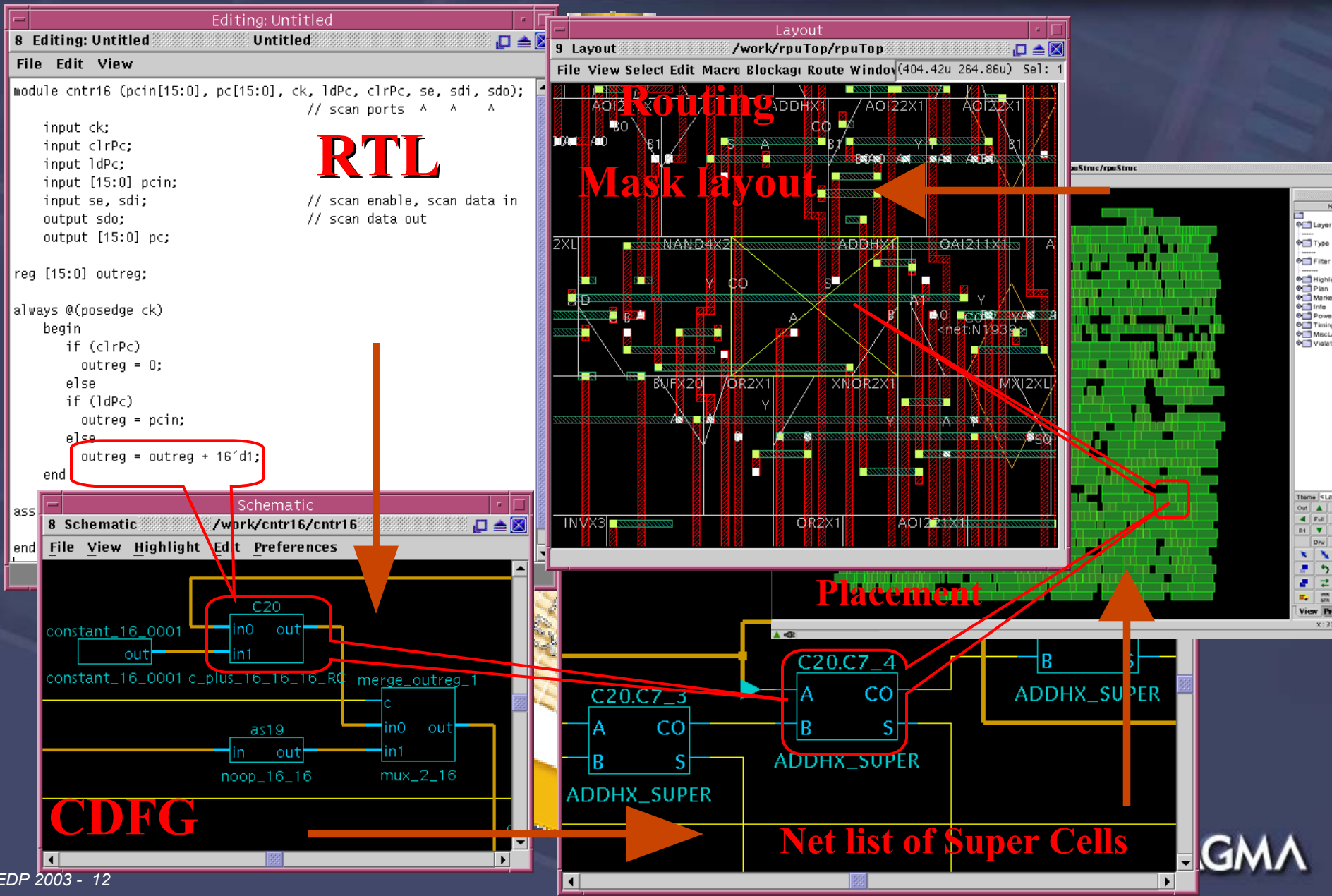
- CPU speeds increasing *much* faster than disk speeds
- Design size growing exponentially
- Disk I/O dominates runtime!
- RAM is very cheap
- 64-bit CPUs enable unlimited data sizes

# Magma's Unified Datamodel



- All tools share a common data structure. They run *directly* on it.
  - Highest speed
  - Lowest memory overhead
- All design data lives “in core”. The tools run *around* the data.
- Data model that contains *everything*:
  - Complete net list, constraints, layout, library data, congestion, etc.
- Analysis tools are part of the core
  - timer, extractor, DRC, power, etc.

# Tight & efficient tool integration



# Example of the strength of the data model: Supporting DSM crosstalk repair



- **Post-route Silicon repair flow:**
  - Update parasitic extraction
  - Update timer
  - Determine/filter problem (e.g. crosstalk noise)
  - Re-route victim:
    - redo-global route
    - redo-detailed route
  - And/or resize and buffer:
    - insert & place gate
    - redo-global route
    - redo-detailed route
  - Update parasitic extraction
  - Update timer
  - Etc. etc. etc.

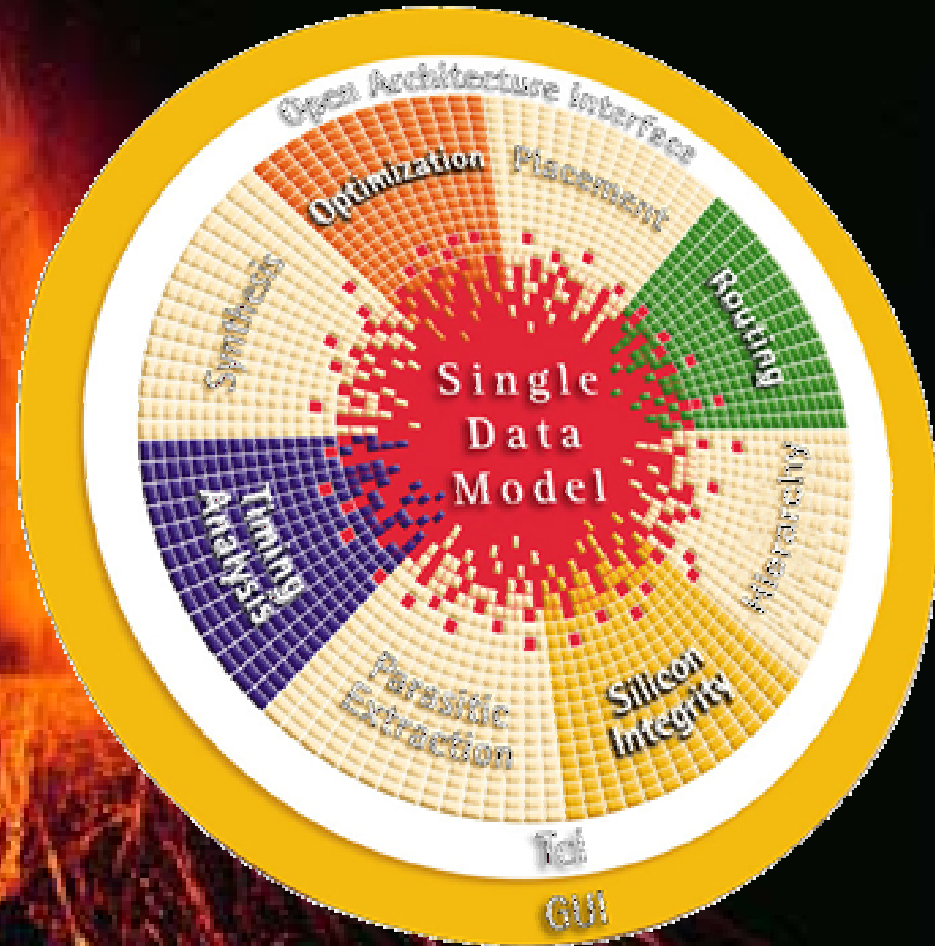
- Similar flows apply to hold time fixing, PVT issues, etc

# Common Data Model supports incremental design and analysis

- **You can add, change or delete**
  - any object in the datamodel
  - at any time
- **Changes are tracked by the data model, it keeps itself consistent at all times**
- **The incremental analysis tools detect changes automatically**
  - ... and update only the affected parts
- **The timer, extractor and DRC engines are brutally incremental**
  - They are never run explicitly
- **Result: all tools have access to most current data**
  - Massive tool interoperability

# Interfacing with the data model through TCL

- **Complete access:**
  - inspect, modify or delete any object or attribute
- **TCL scripts drive the flow**
- **The Graphical User Interface communicates through TCL**
  - easy configuration and adaptation



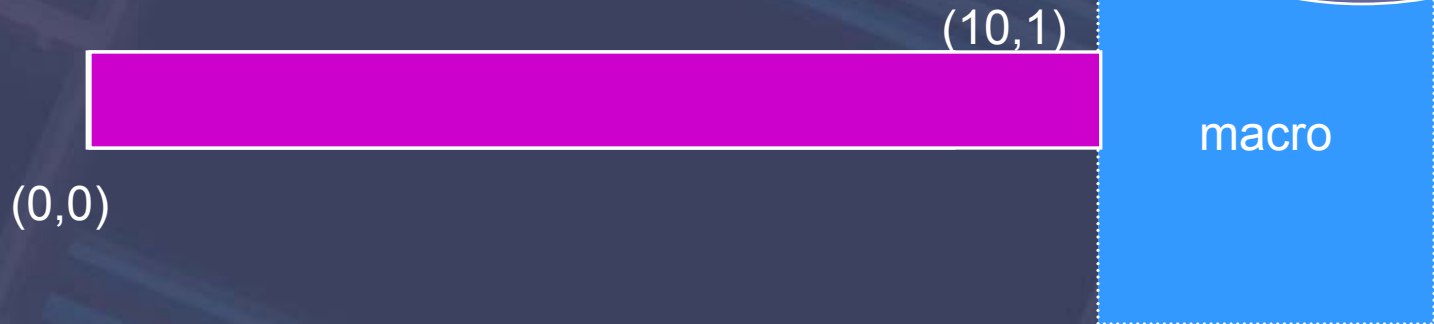
# Example: interaction through TCL

- Create a box (a M1 wire owned by a net called 'newnet')

```
set n [data create net $m -name newnet]
set box "0 0 10u 1u M1 routing $newnet"
data create box $box
```

set up 10u x 1u  
wire at part of

create the wire  
(= box) in the data  
model



- Stretch the wire such that it touches the macro

```
set macrobox [data only model_outline $macro]
data put $box right [query box left $macrobox]
```



# Common Data Model: the enabling technology for RTL to GDSII Design Closure

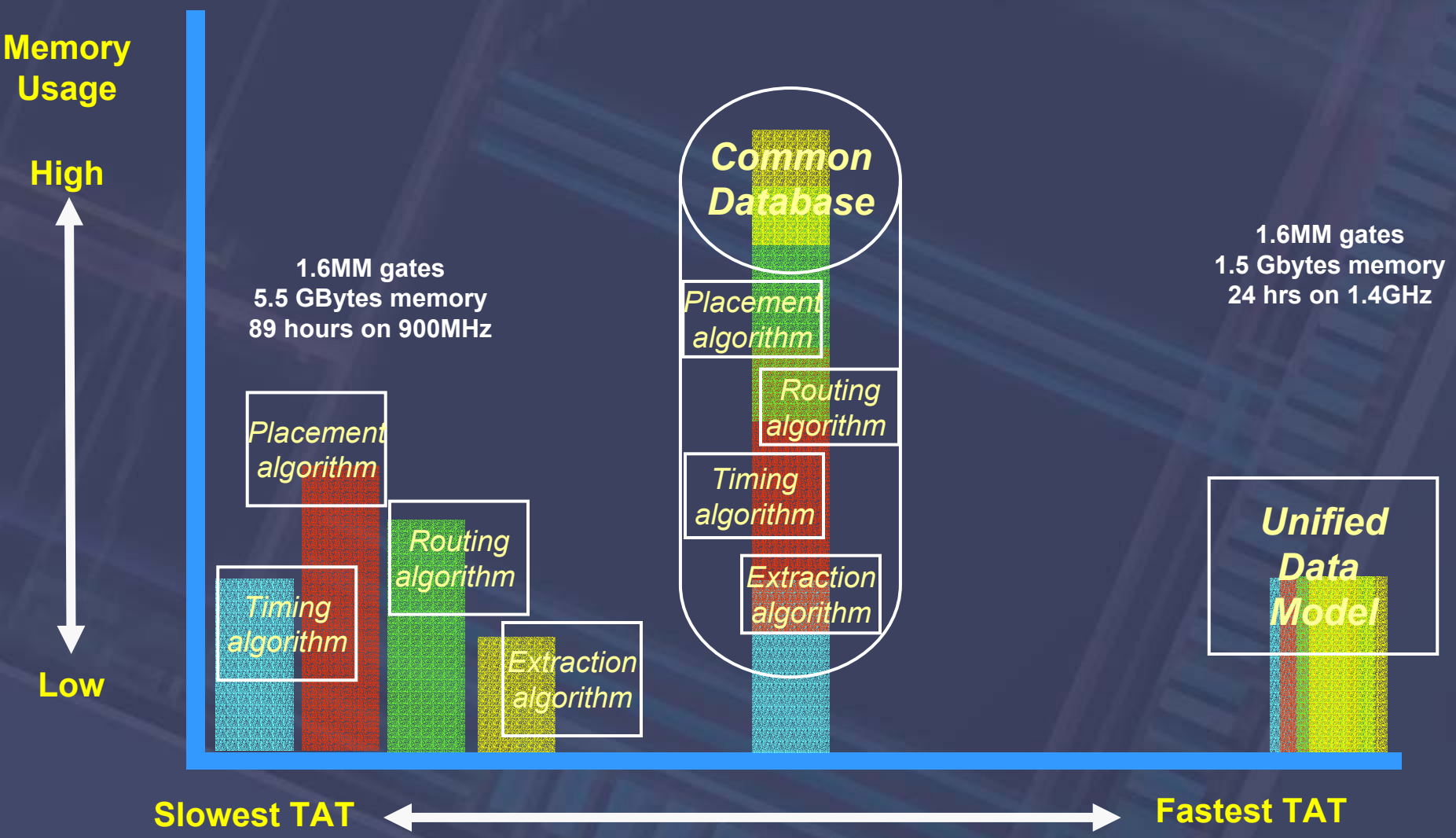
- **Physical Synthesis =**

- Logic optimization
- Placement

- **Design Closure =**

- RTL Synthesis
- Design planning
- Logic optimization
- Placement
- Global routing
- Detailed routing
- Signal Integrity
- Etc...

# Integration through a database will always fall short



# What is OpenAccess?

## The standard

Information Model  
(Graphical)

Data Model  
(C++ headers)

API Specification  
(C++ Binding)

Test Cases  
(C++)

## The reference implementation

Runtime Memory

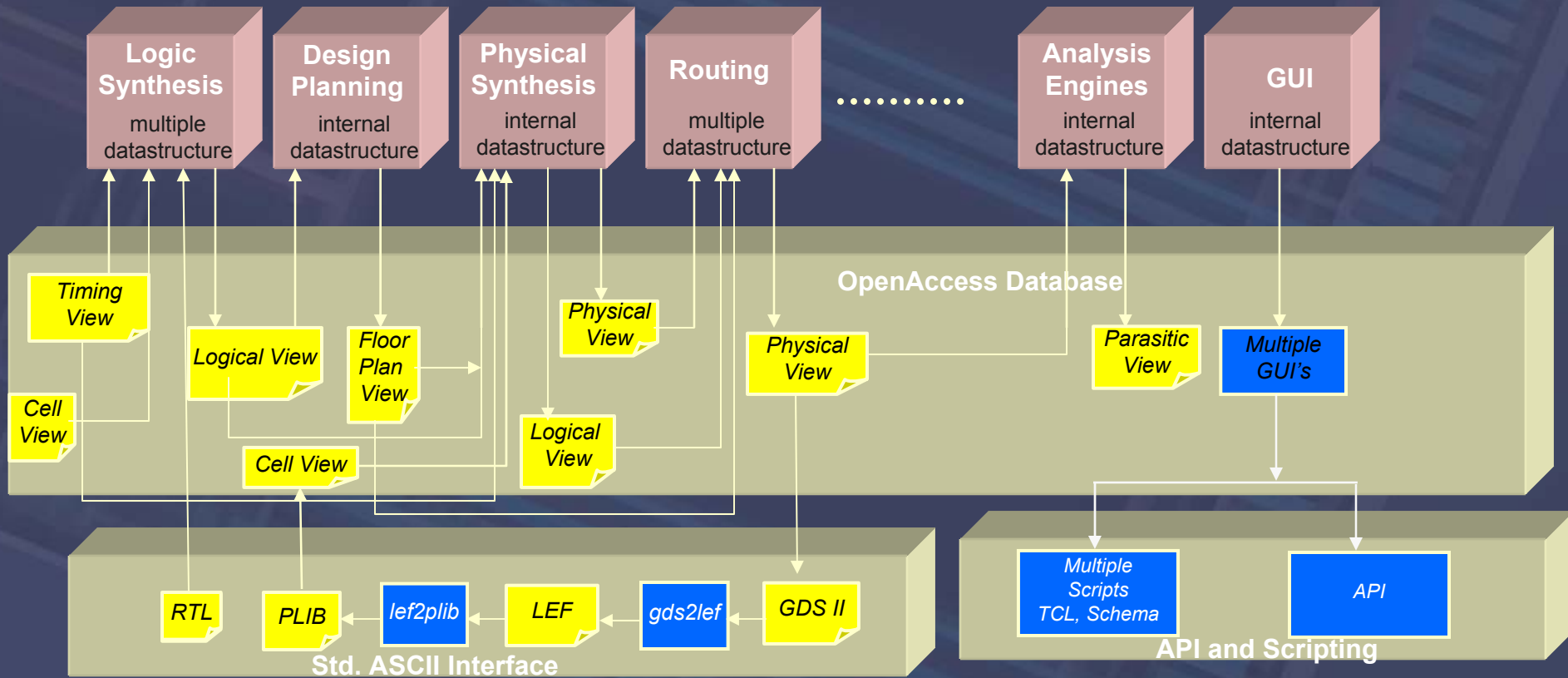
API Implementation  
(solaris/HP-UX, C++ Binding)

Persistent Core

applications

- Open API and source code
- Flexible license model
- High capacity & performance
- Advanced features:
  - Area queryability
  - Callbacks (“don’t abuse these”)
  - Application defined data objects (sparse or dense)
  - Compression/hashing
  - On demand loading
  - Thread safe
  - Access management

# Can OpenAccess be the Common Database?



# Can OpenAccess be the Common Data Model?

- **Yes, in principle, but...**
  - EDA = **F**low + **D**ata model + **A**lgorithms
  - Current generation of point tools not architected for advanced flows
  - Current generation of fully integrated tools are already on the market
    - Extremely wide API to datamodel
  - Proprietary data models provide competitive advantage
    - Control over performance and capacity
    - Carefully tuned to optimize flow
    - Quick turnaround time for bugs and enhancements
  - Many person years already invested in optimizing and debugging proprietary datamodels

# Summary

- **Tight tool integration is crucial  $\leq 90\text{nm}$**
- **Common data model enables integration**
- **“best in class point tools” – nostalgia**
- **Flows, flows, flows**
- **OpenAccess**
  - Magma is committed to interoperability
  - Magma wholeheartedly supports OpenAccess
  - We will read/write to it like any other standard interchange format
  - Unlikely to replace our proprietary data model